Faster ground state preparation and high-precision ground energy estimation with fewer qubits

Yimin Ge, J. Tura, J.I. Cirac QIP 2018

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"Nature isn't classical, dammit, and if you want to make a simulation of nature, you'd better make it quantum mechanical, and by golly it's a wonderful problem, because it doesn't look so easy."



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Quantum simulation



Quantum chemistry





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Small quantum computers





H, t

 $|\psi_0
angle$









Many important applications: $|\psi_0\rangle$ ground state of another non-trivial Hamiltonian!



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Ground state problems generally hard! But may not apply to *natural* systems







QIP 2016



1. General approaches for ground state preparation

2. Algorithms - details

3. Suitability for early quantum computers

Phase estimation











$$H(0) \xrightarrow{H(s)} H(1)$$







Adiabatic algorithms $H(0) \xrightarrow{H(s)} H(1)$

 $|\text{GS}(0)\rangle$











Phase estimation





Phase estimation





Paradigm: First heuristic method, then phase estimation



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This work: Improves part of phase estimation



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Problem: Project given trial state $|\phi\rangle$ onto its ground state component





Setup

 $N \times N$ Hamiltonian H, spectrum in [0, 1]

- Eigenstates $|\lambda_i\rangle$
- Ground energy $\lambda_0,$ ground state $|\lambda_0\rangle$
- All other eigenvalues: $\lambda_i \geq \lambda_0 + \Delta$
- Can efficiently perform time evolution of H

(eg sparse & oracle access, linear combination of easy unitaries, etc [BCK15,BCCKS15,LC16,LC17])



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Circuit \mathcal{C}_{ϕ} , prepares trial state $|\phi
angle$

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- Trivial assumption: $\chi = e^{-O(\log N)}$



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Aim: Extract state $|\lambda'_0\rangle$ st $|| |\lambda'_0\rangle - |\lambda_0\rangle || < \epsilon$ for given ϵ



Results & Comparisons



Results & Comparisons

Ground state preparation



Results & Comparisons

Ground state preparation

Ground energy known


Ground state preparation



- N = total dimension of H
- $\Delta =$ known lower bound on spectral gap
- $\epsilon = \text{allowed error}$

- $\phi_0={\rm overlap}$ of trial state with ground state
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- Λ = base cost of Hamiltonian simulation
- $\Phi = \text{cost of preparing trial state } |\phi\rangle$



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Ground state preparation

Algorithm	Gates	Qubits
$Phase\ est + AA$	$ ilde{O}\left(rac{\Lambda}{ \phi_0 ^2\Delta\epsilon}+rac{\Phi}{ \phi_0 } ight)$	$O\left(\log N + \log rac{1}{\epsilon} + \log rac{1}{\Delta} ight)$
Multicopy PEA (eg [PW'09])	$ ilde{O}\left(rac{\Lambda}{ \phi_0 \Delta}+rac{\Phi}{ \phi_0 } ight)$	

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Ground energy estimation



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Ground energy estimation



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Phase est	$ ilde{O}\left(rac{\Lambda}{\chi^3\xi}+rac{\Phi}{\chi} ight)$	$O\left(\log N + \log \frac{1}{\xi}\right)$
Multicopy PEA (eg [PW'09])	$\tilde{O}\left(\frac{\Lambda}{\chi\xi^{3/2}}+\frac{\Phi}{\chi\sqrt{\xi}}\right)$	$O\left(\log \textit{N} + rac{\lograc{1}{\chi}}{\log\lograc{1}{\chi}} imes \lograc{1}{\xi} ight)$
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- $\epsilon = \mathsf{allowed} \; \mathsf{error}$
- ξ = required precision

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Idea:



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Implementing linear combination of unitaries

eg [CKS'15]

LCU Lemma: Able to perform unitaries $U_k \Rightarrow$ can perform $V := \sum_k \alpha_k U_k$

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$$\begin{array}{c} 3\left|0\right\rangle = \frac{1}{\sqrt{\alpha}}\sum\sqrt{\alpha_{k}}\left|k\right\rangle, \quad \alpha = \sum\left|\alpha_{k}\right.\\ \left|0\right\rangle & \boxed{B} & \boxed{B^{\dagger}} & \left\langle0\right|\\ \left|\phi\right\rangle & \boxed{U_{k}} & V\left|\phi\right\rangle \end{array}$$

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Postselection on ancilla: implement V deterministically

Implementing linear combination of unitaries

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$$|*
angle = rac{1}{lpha} \left|0
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angle \, oldsymbol{V} \left|\phi
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angle + \sqrt{1 - rac{1}{lpha^2}} \left|R
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angle, \qquad \qquad \langle 0|R
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2. Amplitude amplification:

$$\left\| \frac{1}{\alpha} \left| 0 \right\rangle \mathbf{V} \left| \psi \right\rangle \right\| \to 1$$

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2. Approximate as linear combination of easy unitaries

Assume: ground energy known. $H' := H - \lambda_0$



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 $\cos^{2m} H'$

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 $\cos^{2m}H'\ket{\phi} \stackrel{\propto}{_\sim} \ket{\lambda_0}$ for $m pprox 1/\Delta^2$

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$$\cos^{2m} H' = \sum_{k=-m}^{m} \alpha_k e^{-2iH'k}, \quad \alpha_k := \frac{1}{2^{2m}} \binom{2m}{m+k}$$

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3. Use LCU Lemma

Alternative: 1. $(1 - H'^2)^{2m}$ as approximate ground state projector

- 2. Expand in Chebyshev polynomials
- 3. Quantum walks

Algorithm – ground energy unknown



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Previous algorithm:

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• Requires knowing ground energy up to precision $\tilde{O}(\Delta)$



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- Smaller values OK, but exponentially small prob of success

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Lemma (Minimum label finding)

- L unitaries $U_j |0\rangle |0\rangle = |0\rangle |\Phi_j\rangle + |R_j\rangle$, $\langle 0|R_j\rangle = 0$
- $|\Phi\rangle := \frac{1}{\sqrt{L}} \sum_{j} |0\rangle |j\rangle |\Phi_{j}\rangle + |R\rangle$, $\langle 0|R\rangle = 0$

 $\Rightarrow \text{ Given } \chi, \text{ can approximately find smallest } j \text{ s.t. } \||\Phi_j\rangle\| \ge \chi$ using $\tilde{O}(\sqrt{L}/\chi)$ calls to $U = \sum_j |j\rangle\langle j| \otimes U_j$
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Idea: Binary search on label ancilla using amplitude amplification

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- $U_j =$ previous algorithm, assuming ground energy is $E_j \gtrsim j\Delta$
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- Runtime dependence on χ , not $|\phi_0|$

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Bonus: Also find ground energy to precision $\tilde{O}(\Delta)$

 Δ only required to be lower bound on gap

 $\Rightarrow~$ general ground energy estimation algorithm for high precisions

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 $\Rightarrow~$ general ground energy estimation algorithm for high precisions

Alternative: first use PEA to find ground energy

 $\rightarrow~$ better scaling in Δ but worse scaling in overlap



Adaption for early quantum computers:



Adaption for early quantum computers: Amplitude amplification



Adaption for early quantum computers: <u>Amplitude amplification</u> Repeated measurements



Adaption for early quantum computers: Amplitude amplification Repeated measurements

NISQ: devices with $\approx 100~\text{qubits}$



Adaption for early quantum computers: <u>Amplitude amplification</u> Repeated measurements



NISQ: devices with ≈ 100 qubits, $\approx 10^4 - 10^5$ (?) gates reliably

Limiting factor: number of gates coherently in *single-run*, **not** *total* runtime!

Adaption for early quantum computers: Amplitude amplification Repeated measurements



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Ground state preparation algorithms, ground energy known

Algorithm

Multicopy PEA

Phase est

This work

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 ϵ = allowed error

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Alg	orithm			
Multicopy PEA Too many		Too many o	qubits	
Pha	ase est			
Thi	is work			
	N = total dimension of $H\Delta = known lower bound on spectral gap\epsilon = allowed error$		ϕ_0 = overlap of trial state with ground state Λ = base cost of Hamiltonian simulation Φ = cost of preparing trial state $ \phi\rangle$	

Adaption for early quantum computers: <u>Amplitude amplification</u> Repeated measurements



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Alac	rithm	Gates		
Aigc	gorithm	Ampl amplif		
Multicopy PEA		EA	Too many qubits	
Pha	se est	$ ilde{O}\left(rac{\Lambda}{ \phi_0 ^2\Delta\epsilon}+rac{\Phi}{ \phi_0 } ight)$		
This	s work	$ ilde{O}\left(rac{\Lambda}{ \phi_0 \Delta}+rac{\Phi}{ \phi_0 } ight)$		
	N = total dimension of H		$\phi_{\rm O}={\rm overlap}$ of trial state with ground state	
$\Delta = know$		lower bound on spectral gap	$\Lambda=$ base cost of Hamiltonian simulation	
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Algorithm	Gates	Gates
Algorithm	Ampl amplif	Repeated mmt
Multicopy PEA		Too many qubits
Phase est	$ ilde{O}\left(rac{\Lambda}{ \phi_0 ^2\Delta\epsilon}+rac{\Phi}{ \phi_0 } ight)$	$ ilde{O}\left(rac{\Lambda}{ \phi_0 ^3\Delta\epsilon}+rac{\Phi}{ \phi_0 ^2} ight)$
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Almonithms	Gates	Gates	Gates	Depatition	
Algorithm	Ampl amplif	Repeated mmt	single run	Repetitions	
Multicopy PEA		Too many qubits			
Phase est	$\tilde{O}\left(\frac{\Lambda}{ \phi_0 ^2\Delta\epsilon}+\frac{\Phi}{ \phi_0 } ight)$	$\tilde{O}\left(rac{\Lambda}{ \phi_0 ^3\Delta\epsilon}+rac{\Phi}{ \phi_0 ^2} ight)$	$\tilde{O}\left(rac{\Lambda}{ \phi_0 \Delta\epsilon}+\Phi ight)$	$\tilde{O}\left(rac{1}{ \phi_0 ^2} ight)$	
This work	$ ilde{O}\left(rac{\Lambda}{ \phi_0 \Delta}+rac{\Phi}{ \phi_0 } ight)$	$ ilde{O}\left(rac{\Lambda}{ \phi_0 ^2\Delta}+rac{\Phi}{ \phi_0 ^2} ight)$	$ ilde{O}\left(rac{\Lambda}{\Delta}+\Phi ight)$	$ ilde{O}\left(rac{1}{ \phi_0 ^2} ight)$	
N = total	dimension of H	$\phi_0 = or$	verlap of trial state with g	round state	
$\Delta = know$	vn lower bound on spectral gap	$\Lambda = bas$	se cost of Hamiltonian sim	ulation	
$\epsilon = \text{allowe}$	ed error	$\Phi = \cos \theta$	st of preparing trial state	$ \phi\rangle$	

Adaption for early quantum computers: <u>Amplitude amplification</u> Repeated measurements

NISQ: devices with ≈ 100 qubits, $\approx 10^4 - 10^5$ (?) gates reliably Limiting factor: number of gates coherently in *single-run*, **not** *total* runtime!

Algorithm	Gates Ampl amplif	Gates Repeated mmt	Gates single run	Repetitions
Phase est	$ ilde{O}\left(rac{\Lambda}{ \phi_0 ^2\Delta\epsilon}+rac{\Phi}{ \phi_0 } ight)$	$\tilde{O}\left(rac{\Lambda}{ \phi_0 ^3\Delta\epsilon}+rac{\Phi}{ \phi_0 ^2} ight)$	$ ilde{O}\left(rac{\Lambda}{ \phi_0 \Delta\epsilon}+\Phi ight)$	$\tilde{O}\left(rac{1}{ \phi_0 ^2} ight)$
This work	$ ilde{O}\left(rac{\Lambda}{ \phi_0 \Delta}+rac{\Phi}{ \phi_0 } ight)$	$\tilde{O}\left(\frac{\Lambda}{ \phi_0 ^2\Delta}+\frac{\Phi}{ \phi_0 ^2}\right)$	$ ilde{O}\left(rac{\Lambda}{\Delta}+\Phi ight)$	$\tilde{O}\left(rac{1}{ \phi_0 ^2} ight)$

- N = total dimension of H
- $\Delta = \mathsf{known} \ \mathsf{lower} \ \mathsf{bound} \ \mathsf{on} \ \mathsf{spectral} \ \mathsf{gap}$
- $\epsilon = \mathsf{allowed} \; \mathsf{error}$

- $\phi_0={\rm overlap}$ of trial state with ground state
- $\Lambda=$ base cost of Hamiltonian simulation
- $\Phi=\mathrm{cost}$ of preparing trial state $|\phi\rangle$



Ground state preparation algorithm

- Faster than naive phase estimation
- Fewer qubits than improved phase estimation
- Known and unknown ground energy
- Estimates ground energy to high precision



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Applications

- Quantum simulation of many-body systems (quenches!)
- Quantum chemistry
- Single-copy tomography, QMA witnesses, optimisation problems, quantum machine learning, ...

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Potential applications for early quantum computers!